# **Consensus on Demand**

Jakub Sliwinski, Yann Vonlanthen, Roger Wattenhofer

ETH Zurich - **Dis**tributed **Co**mputing Group

#### Bitcoin: A Peer-to-Peer Electronic Cash System

Satoshi Nakamoto satoshin@gmx.com www.bitcoin.org

Abstract. A purely peer-to-peer version of electronic cash would allow online payments to be sent directly from one party to another without going through a financial institution. Digital signatures provide part of the solution, but the main benefits are lost if a trusted third party is still required to prevent double-spending. We propose a solution to the double-spending problem using a peer-to-peer network. The network timestamps transactions by hashing them into an ongoing chain of hash-based proof-of-work, forming a record that cannot be changed without redoing the proof-of-work. The longest chain not only serves as proof of the sequence of events witnessed, but proof that it came from the largest pool of CPU power. As long as a majority of CPU power is controlled by nodes that are not cooperating to attack the network, they'll generate the longest chain and outpace attackers. The network itself requires minimal structure. Messages are broadcast on a best effort basis, and nodes can leave and rejoin the network at will, accepting the longest proof-of-work chain as proof of what happened while they were gone.





## **Prevent Double-Spending**



## **Prevent Double-Spending**



# **Prevent Double-Spending**

Total Order



# **Prevent Double-Spending**

## Total Order = Consensus



Prevent Double-Spending f Total Order = Consensus



Prevent Double-Spending f Total Order = Consensus

# Online Payments Without Consensus

A Non-Consensus Based Decentralized Financial Transaction Processing Model with Support for Efficient Auditing

by

Saurabh Gupta

A Thesis Presented in Partial Fulfillment of the Requirements for the Degree Master of Science

# Online Payments Without Consensus

A Non-Consensus Based Decentralized Financial Transaction Processing Model with Support for Efficient Auditing

by

Saurabh Gupta

A Thesis Presented in Partial Fulfillment of the Requirements for the Degree Master of Science

#### ABC: Asynchronous Blockchain without Consensus

Jakub Sliwinski and Roger Wattenhofer

ETH Zurich {jsliwinski,wattenhofer}@ethz.ch

Abstract. There is a preconception that a blockchain needs consensus. But consensus is a powerful distributed property with a remarkably high price tag. So one may wonder whether consensus is at all needed. We introduce a new blockchain architecture called ABC that functions despite not establishing consensus, and comes with an array of advantages: ABC is permissionless, deterministic, and resilient to complete asynchrony. ABC features finality and does not rely on costly proof-ofwork.

Without establishing consensus, ABC cannot support certain applica-

# Online Payments Without Consensus

# Online Payments Without Consensus

A Non-Consensus Based Decentralized Financial Transaction Processing Model with Support for Efficient Auditing

by

Saurabh Gupta

#### The Consensus Number of a Cryptocurrency

Rachid Guerraoui rachid.guerraoui@epfl.ch EPFL Lausanne. Switzerland

Petr Kuznetsov petr.kuznetsov@telecom-paristech.fr LTCI, Télécom Paris, IP Paris Paris, France Matteo Monti matteo.monti@epfl.ch EPFL Lausanne, Switzerland

distributed computing, distributed asset transfer, blockchain, con-

Rachid Guerraoui, Petr Kuznetsov, Matteo Monti, Matej Pavlovič, and Dragos-Adrian Seredinschi. 2019. The Consensus Number of a Cryptocurrency. In

2019 ACM Symposium on Principles of Distributed Computing (PODC'19).

July 29-August 2, 2019, Toronto, ON, Canada. ACM, New York, NY, USA,

Matej Pavlovič matej.pavlovic@epfl.ch EPFL Lausanne. Switzerland Dragos-Adrian Seredinschi\* dragos-adrian.seredinschi@epfl.ch EPFL Lausanne. Switzerland

10 pages. https://doi.org/10.1145/3293611.3331589

KEYWORDS

ACM Reference Format

#### ABSTRACT

Many blockchain-based algorithms, such as Bitcoin, implement a desentialized asstranafer system, often referret to as a cryptourrrncy. As stated in the original paper by Nakamoto, at the heat of these systems lists the problem of preventing double-gending, this is usually solved by achieving consensus on the order of transfers mong the participants. By treating the asset transfer problem as a concurrent object and determining its consensus number, we show that consensus is not necessary to prevent double-gending. We first consider the problem as defined by Nakamoto, where only a sindle process—the account oncurre-can withdhard from each

only a single process-the account owner-can withdraw from ear

Jakub Sliwinski and Roger Wattenhofer

ETH Zurich {jsliwinski,wattenhofer}@ethz.ch

Abstract. There is a preconception that a blockchain needs consensus. But consensus is a powerful distributed property with a remarkably high price tag. So one may wonder whether consensus is at all needed. We introduce a new blockchain architecture called ABC that functions despite not establishing consensus, and comes with an array of advantages: ABC is permissionless, deterministic, and resilient to complete asynchrony. ABC features finality and does not rely on costly proof-ofwork.

Without establishing consensus, ABC cannot support certain applica-



























# What benefits do we have?

- Low latency
- Parallel execution →
   Horizontal scaling



















# What can we do with this system?











Only done when necessary!

# Consensus on Demand

#### Consensus on Demand

Jakub Sliwinski, Yann Vonlanthen, and Roger Wattenhofer

ETH Zurich, Switzerland

Abstract. Digital money can be implemented efficiently by avoiding consensus. However, no-consensus implementations have drawbacks, as they cannot support smart contracts, and (even more fundamentally) they cannot deal with conflicting transactions.

We present a novel protocol that combines the benefits of an asynchronous, broadcast-based digital currency, with the capacity to perform consensus. This is achieved by selectively performing consensus a posteriori, i.e., only when absolutely necessary. Our on-demand consensus comes at the price of restricting the Byzantine participants to be less than a one-fifth minority in the system, which is the optimal threshold. We formally prove the correctness of our system and present an opensource implementation, which inherits many features from the Ethereum ecosystem.

Keywords: Blockchain  $\cdot$  Byzantine fault tolerance  $\cdot$  Consensus  $\cdot$  Cryptocurrencies  $\cdot$  Reliable broadcast.



















# The Best of Both Worlds

#### Fast Path

- Low latency
- Parallelization
- + Lower message complexity
  - No need for BRB!

#### Consensus

- Arbitrary Turing complete computation
  - Enables smart contracts





## Why do we need $n \ge 5f + 1$ ?

# 

## Let's assume n = 5f







# Slow path must agree

with fast path!



# How can we guarantee agreement?







In case of any conflict, propose most observed *transaction* after (n+3f)/2 observed *acks*.





# Implementation







User

|                           | Bitcoin and<br>Ethereum [32] | Ouroboros<br>[24] | Algorand [19] | PBFT [12]             | Red Belly [15]        | BEAT [17]    | Broadcast-<br>based [14] | $C_{oD}+PBFT$ | CoD + Br. |
|---------------------------|------------------------------|-------------------|---------------|-----------------------|-----------------------|--------------|--------------------------|---------------|-----------|
| Energy-efficient          |                              | ~                 | $\checkmark$  | <ul> <li>✓</li> </ul> | <ul> <li>✓</li> </ul> | $\checkmark$ | ✓                        | $\checkmark$  | / v       |
| Deterministic<br>finality |                              |                   | $\checkmark$  | $\checkmark$          | $\checkmark$          | $\checkmark$ | $\checkmark$             | √             | ~         |
| Permissionless            | $\checkmark$                 | $\checkmark$      | $\checkmark$  |                       |                       |              | 6°                       |               |           |
| Leaderless                |                              |                   |               |                       | $\checkmark$          | $\checkmark$ | $\checkmark$             |               | ~         |
| Asynchronous              |                              |                   |               |                       |                       | $\checkmark$ | $\checkmark$             |               | ~         |
| Parallelizable            |                              |                   |               |                       |                       |              | $\checkmark$             | $\checkmark$  | ~         |
| Consensus                 | $\checkmark$                 | $\checkmark$      | $\checkmark$  | $\checkmark$          | ~                     | $\checkmark$ |                          | $\checkmark$  | ~         |

# We provide a wrapper that minimizes the accesses to Consensus!

# Thank you! Merci!

### yvonlanthen@ethz.ch

ETH Zurich - Distributed Computing Group