Eingebettete Systeme
Echtzeitverhalten und Betriebssysteme

2. Grundlegende Modelle

Basic terms

- **Real-time systems**
  - **Hard**: A real-time task is said to be hard, if missing its deadline may cause catastrophic consequences on the environment under control. Examples are sensory data acquisition, detection of critical conditions, actuator servoing.
  - **Soft**: A real-time task is called soft, if meeting its deadline is desirable for performance reasons, but missing its deadline does not cause serious damage to the environment and does not jeopardize correct system behavior. Examples are command interpreter of the user interface, displaying messages on the screen.

Schedule

- **Given a set of tasks** \( J = \{J_1, J_2, \ldots\} \):
  - A **schedule** is an assignment of tasks to the processor, such that each task is executed until completion.
  - A schedule can be defined as an integer step function \( \sigma : R \to N \) where \( \sigma(t) \) denotes the task which is executed at time \( t \). If \( \sigma(t) = 0 \) then the processor is called **idle**.
  - If \( \sigma(t) \) changes its value at some time, then the processor performs a **context switch**.
  - Each interval, in which \( \sigma(t) \) is constant is called a **time slice**.
  - A **preemptive schedule** is a schedule in which the running task can be arbitrarily suspended at any time, to assign the CPU to another task according to a predefined scheduling policy.
Schedule and timing

- A schedule is said to be **feasible**, if all task can be completed according to a set of specified constraints.
- A set of tasks is said to be **schedulable**, if there exists at least one algorithm that can produce a feasible schedule.
- **Arrival time** $a_i$, or **request or release time** $r_i$, is the time at which a task becomes ready for execution.
- **Computation time** $C_i$ is the time necessary to the processor for executing the task without interruption.
- **Deadline** $d_i$ is the time at which a task should be completed.
- **Start time** $s_i$ is the time at which a task starts its execution.
- **Finishing time** $f_i$ is the time at which a task finishes its execution.

Using the above definitions, we have $d_i \geq r_i + C_i$

**Lateness** $L_i = f_i - d_i$ represents the delay of a task completion with respect to its deadline; note that if a task completes before the deadline, its lateness is negative.

**Tardiness** or **exceeding time** $E_i = \max(0, L_i)$ is the time a task stays active after its deadline.

**Laxity or slack time** $X_i = d_i - a_i - C_i$ is the maximum time a task can be delayed on its activation to complete within its deadline.

**Example**

Computation times: $C_1 = 9, C_2 = 12$
Start times: $s_1 = 0, s_2 = 6$
Finishing times: $f_1 = 18, f_2 = 28$
Lateness: $L_1 = -4, L_2 = 1$
Tardiness: $E_1 = 0, E_2 = 1$
Laxity: $X_1 = 13, X_2 = 11$
Precedence constraints

- Precedence relations between graphs can be described through an acyclic directed graph $G$ where tasks are represented by nodes and precedence relations by arrows. $G$ induces a partial order on the task set.

- There are different interpretations possible:
  - All successors of a task are activated (concurrent task execution).
  - One successor of a task is activated (non-deterministic choice).

Example (concurrent activation):
- Image acquisition $acq_1$ $acq_2$
- Low level image processing $edge_1$ $edge_2$
- Feature/contour extraction $shape$
- Pixel disparities $disp$
- Object size $H$
- Object recognition $rec$

Resource constraints

- Examples of common resources: data structures, variables, main memory area, file, set of registers, I/O unit, ... .

- Many shared resources do not allow simultaneous accesses but require mutual exclusion (exclusive resources). A piece of code executed under mutual exclusion constraints is called a critical section.

Example: $J_1$ has higher priority than $J_2$

Process state diagram:
Classification of scheduling algorithms

- **With preemptive algorithms**, the running task can be interrupted at any time to assign the processor to another active task, according to a predefined scheduling policy.
- **With a non-preemptive algorithm**, a task, once started, is executed by the processor until completion.
- **Static algorithms** are those in which scheduling decisions are based on fixed parameters, assigned to tasks before their activation.
- **Dynamic algorithms** are those in which scheduling decisions are based on dynamic parameters that may change during system execution.

Classification of scheduling algorithms

- **An off-line algorithm** is executed on the entire task set before the actual task activation. The schedule generated in this way is stored in a table and later executed by a dispatcher. Advantage: Small overhead even in case of complex scheduling algorithms.
- **An on-line algorithm** takes scheduling decisions at runtime every time a new task enters the system or when a running task terminates. Advantage: Can adapt to dynamic task sets.
- **Guarantee-based scheduling** algorithms apply an acceptance test to each new task $J_{new}$. If it cannot be planned together with already granted tasks, it is rejected; otherwise it is accepted and added to the current set of tasks: $J_{old} \cup \{J_{new}\}$.

Classification of scheduling algorithms

A new task $J_{new}$ is accepted and added to the current set of tasks: $J_{old} \cup \{J_{new}\}$.

Scheduling methods and metrics

- **Stochastic models**: Queuing theory, Markov chains give no deterministic guarantees. Tasks meet their deadlines only with a certain probability.
- **Deterministic models**: Worst-case analysis gives guarantees but may be overly pessimistic in most of the instances.
Metrics

- **Average response time:**
  \[ \bar{r} = \frac{1}{n} \sum_{i=1}^{n} (f_i - r_i) \]
- **Total completion time:**
  \[ t_c = \max_i (f_i - \min_i(r_i)) \]
- **Weighted sum of completion time:**
  \[ t_w = \frac{\sum_{i=1}^{n} w_i(f_i - \eta_i)}{\sum_{i=1}^{n} w_i} \]
- **Maximum lateness:**
  \[ L_{\text{max}} = \max_i (f_i - d_i) \]
- **Maximum number of late tasks:**
  \[ N_{\text{late}} = \sum_i \text{miss}(f_i) \]
  \[ \text{miss}(f_i) = \begin{cases} 0 & \text{if } f_i \leq d_i \\ 1 & \text{otherwise} \end{cases} \]

Metrics example

![Diagram showing two tasks with different response times and completion times.](image)

**Average response time:**
\[ \bar{r} = \frac{1}{2} (18 + 24) = 21 \]
**Total completion time:**
\[ t_c = 28 - 0 = 28 \]
**Weighted sum of completion time:**
\[ t_w = \frac{2 + 18 + 24}{3} = 20 \]
**Maximum number of late tasks:**
\[ N_{\text{late}} = 1 \]
**Maximum lateness:**
\[ L_{\text{max}} = 1 \]

Scheduling example

- In (a), the maximum lateness is minimized, but all tasks miss their deadlines.
- In (b), the maximal lateness is larger, but only one task misses its deadline.